

High Performance Traffic Shaping for DDoS Mitigation

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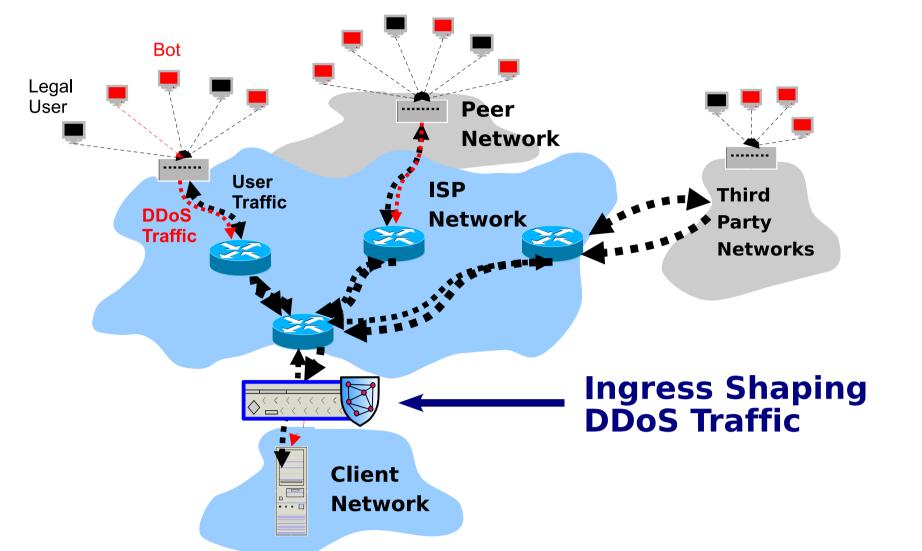
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Introduction

- DDoS flooding attacks (UDP, ICMP, SYN) easy identifiable
- TCP based protocol conformant DDoS attacks (e.g. HTTP-GET floods) have lower deviation compared to normal user traffic
- \rightarrow more difficult to distinguish valid and illegal traffic
- \rightarrow firewall DROP rules may cause blocking legal users

Idea

• Instead of blocking/accepting flows we shape their source IPs



Shaping Algorithm

- 1: function PACKET_HANDLER(Packet p)
- $r \leftarrow range including. p.source_IP using binary search$
- 3: **if** r not found **then**
 - accept(p) and return
- 5: $q \leftarrow \mathsf{queue of } r$

4:

7:

- 6: **if** not q.empty or r.sent+p.size > r.limit **then**
 - if q.size < q.max_size then
- q.push(p)
 - steel(p)
 - else drop(p)



• Conditional Legitimate Probability (CLP) [5] is the probability of a flow to be legal

$$CLP(p) = \frac{N_n \cdot P_n(A = a_p) \cdot P_n(B = b_n) \cdot \dots}{N_m \cdot P_m(A = a_p) \cdot P_m(B = b_n) \cdot \dots}$$
(1)

- where N_n is the number of normal packets, N_m the number of measured packets (mixture of normal and attack traffic) and $P(A = a_p)$ the probability of a feature A to be a_p .
- CLP is calculated on the basis of previously observed traffic, e.g. histograms on source IP prefixes, packet sizes or server ports

• The higher the CLP, the higher the assigned bandwidth during a DDoS incident **Challenge**

- During an DDoS attack, thousands of source IP addresses need to be shaped.
- On Linux systems, tc [2] can only handle few shaping rules \rightarrow Need for a better performing algorithm for the DDoS scenario.

11:	else
12:	r.sent + = p.size
13:	accept(p)
14: f	unction TIMER_HANDLER
15:	for all ranges r do
16:	$r.sent \leftarrow 0; finished \leftarrow false$
17:	$q \leftarrow queue \ of \ r$
18:	while not q.empty and not finished do
19:	$p \leftarrow q.front()$
20:	<pre>if r.sent + p.size < r.limit then</pre>
21:	send(p)
22:	q.pop()
23:	r.sent += p.size
24:	else finished \leftarrow true
1	

Evaluation

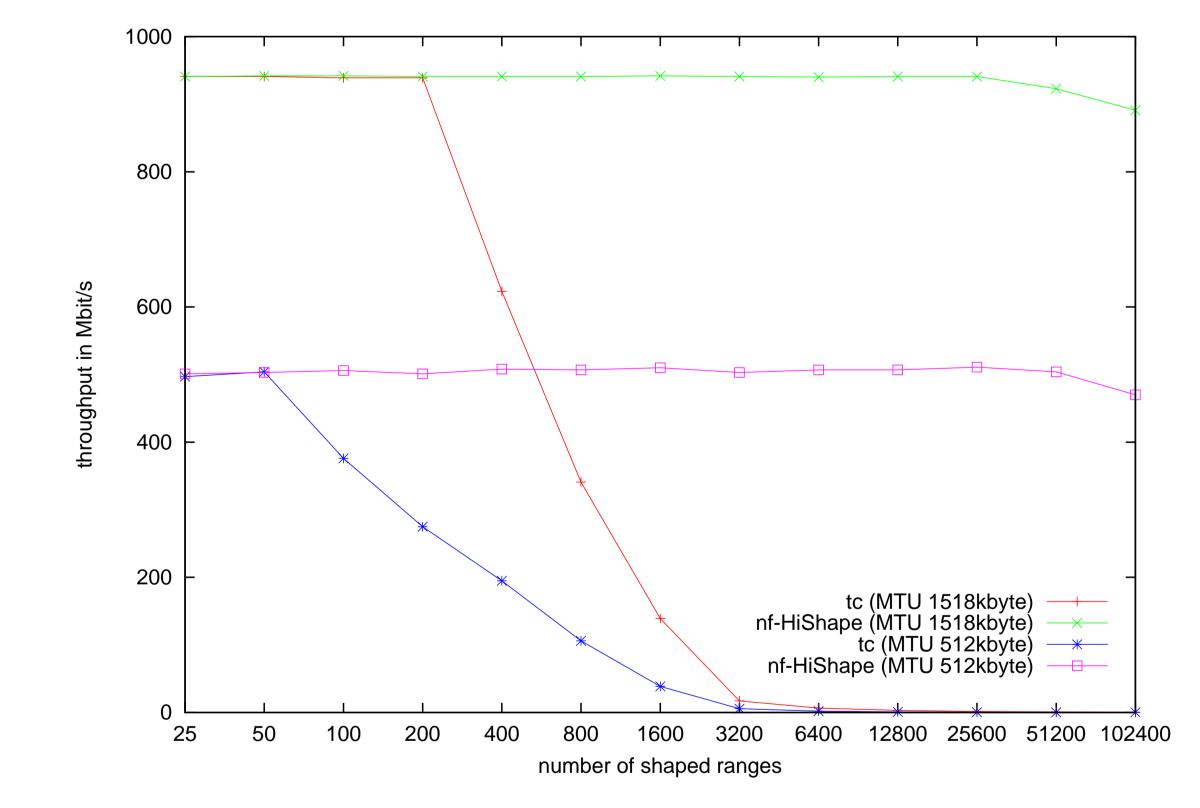
- Measuring throughput of a legal (not shaped) user on a 1GBit/s link depending on the number of shaped IP ranges.
- tc throughput decreases at 400 shaped ranges, not enough to mitigate DDoS attacks
- *nf-HiShape* still performant at 100,000 defined IP ranges

Traffic Shaping

- Easy and fast algorithm for high packet rates
- Specialized for DDoS mitigation (filter parameter only source IP)
- \bullet IP ranges are continuous intervals r over IP addresses $[r^{start}, r^{end}]$ with a defined bandwidth limit
- List of ranges can be sorted to perform binary search:

$$\forall r_i, r_j \in R, i < j : r_i^{end} < r_j^{start}$$
(2)

- Every arriving packet is accepted, queued or dropped, similar to *Random Early Detection (RED)* [4] in the PACKET_HANDLER function.
- A triggered function TIMER_HANDLER sends packets (with respect to the defined bandwidth) and calculates the used bandwidth
- Complexity for each incoming packet is $O(\log_2 n)$, where n is the number of IP ranges with bandwidth limits.
- Worst case complexity: every source IP address has a different target bandwidth $O(\log_2 2^{32}) = 32$ lookups for each incoming packet.



Open Source Software

The presented algorithm is implemented as a Linux kernel module called *nf-HiShape* [1].

It is available under GPL license at http://code.google.com/p/nf-hishape/. It ships with a userland tool which syntax is similar to iptables, e.g.:

nf-hishape -i eth0 -f 192.168.0.0 -t 192.168.255.255-15.0

Future Work

- Support more than one parameter for shaping
- Investigate FIS trees [3] with $O(\log_2 \log_2 N)$ lookup complexity

sets a limit for 192.168.0.0/16 to 5.0 kB/s and

nf-hishape -i eth0 -L

lists all shaping ranges for the given interface. The userland tool also reads in ASCII files with bandwidth limits for faster integration with DDoS mitigation systems.

References

- [1] nf-hishape. http://code.google.com/p/nf-hishape/.
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- [3] A. Feldmann and S. Muthukrishnan. Tradeoffs for packet classification. In *INFOCOM*, pages 1193-1202, 2000.
- [4] S. Floyd and V. Jacobson. Random early detection gateways for congestion avoidance. *IEEE/ACM Transactions on Networking*, 1:397-413, 1993.
- [5] Y. Kim, W. C. Lau, M. C. Chuah, and H. J. Chao. Packetscore: A statistics-based packet filtering scheme against distributed denial-of-service attacks *IEEE Transactions on Dependable and Secure Computing*, 03(2):141-155, 2006.
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